ADVANCED DATABASE SYSTEMS

Index Locking & Latching

@Andy_Pavlo // 15-721 // Spring 2018
TODAY'S AGENDA

Index Locks vs. Latches
Latch Implementations
Index Latching (Logical)
Index Locking (Physical)
DATABASE INDEX

A data structure that improves the speed of data retrieval operations on a table at the cost of additional writes and storage space.

Indexes are used to quickly locate data without having to search every row in a table every time a table is accessed.
DATA STRUCTURES

Order Preserving Indexes
→ A tree-like structure that maintains keys in some sorted order.
→ Supports all possible predicates with $O(\log n)$ searches.

Hashing Indexes
→ An associative array that maps a hash of the key to a particular record.
→ Only supports equality predicates with $O(1)$ searches.
The original B-tree from 1972 stored keys + values in all nodes in the tree.
→ More memory efficient since each key only appears once in the tree.

A B+tree only stores values in leaf nodes. Inner nodes only guide the search process.
→ Easier to manage concurrent index access when the values are only in the leaf nodes.
We already know how to use locks to protect objects in the database.

But we have to treat indexes differently because the physical structure can change as long as the logical contents are consistent.
SIMPLE EXAMPLE

A

| 20 | 22 |
SIMPLE EXAMPLE

**Txn #1:** Read ‘22’
SIMPLE EXAMPLE

Txn #1: Read ‘22’
Txn #2: Insert ‘21’
**SIMPLE EXAMPLE**

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Txn #1: Read ‘22’

Txn #2: Insert ‘21’

Txn #1: Read ‘22’
LOCKS VS. LATCHES

Locks
→ Protects the index’s logical contents from other txns.
→ Held for txn duration.
→ Need to be able to rollback changes.

Latches
→ Protects the critical sections of the index’s internal data structure from other threads.
→ Held for operation duration.
→ Do not need to be able to rollback changes.
LOCK-FREE INDEXES

Possibility #1: No Locks
→ Txns don’t acquire locks to access/modify database.
→ Still have to use latches to install updates.

Possibility #2: No Latches
→ Swap pointers using atomic updates to install changes.
→ Still have to use locks to validate txns.
LATCH IMPLEMENTATIONS

Blocking OS Mutex
Test-and-Set Spinlock
Queue-based Spinlock
Reader-Writer Locks

Source: Anastasia Ailamaki
COMPARE- AND- SWAP

Atomic instruction that compares contents of a memory location \( M \) to a given value \( V \)

→ If values are equal, installs new given value \( V' \) in \( M \)
→ Otherwise operation fails

\[
\text{__sync_bool_compare_and_swap}(&M, 20, 30)
\]
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LATCH IMPLEMENTATIONS

Choice #1: Blocking OS Mutex
→ Simple to use
→ Non-scalable (about 25ns per lock/unlock invocation)
→ Example: `std::mutex`
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→ Example: `std::mutex`

```
std::mutex m;
:
m.lock();
// Do something special...
m.unlock();
```
LATCH IMPLEMENTATIONS

Choice #2: Test-and-Set Spinlock (TAS)
→ Very efficient (single instruction to lock/unlock)
→ Non-scalable, not cache friendly
→ Example: `std::atomic<T>`
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→ Very efficient (single instruction to lock/unlock)
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→ Example: `std::atomic<T>`

```cpp
std::atomic_flag latch;
:
while (latch.test_and_set(...)) {
    // Yield? Abort? Retry?
}
```
Choice #2: Test-and-Set Spinlock (TAS)
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    // Yield? Abort? Retry?
}
\end{verbatim}
LATCHE IMPLEMENTATIONS

Choice #3: Queue-based Spinlock (MCS)
→ More efficient than mutex, better cache locality
→ Non-trivial memory management
→ Example: `std::atomic<Latch*>`
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Base Latch

Mellor-Crummey and Scott
LATCH IMPLEMENTATIONS

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Mellor-Crummey and Scott

Base Latch

CPU1
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Base Latch ───────────────────────────────────────────────────────────────────────────────────────────────────────────────────────────
next

CPU1 Latch ───────────────────────────────────────────────────────────────────────────────────────────────────────────────────────────
next

Mellor-Crummey and Scott
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Mellor-Crummey and Scott

Base Latch → CPU1 Latch → CPU2 Latch

next

CPU1

CPU2

CPU3
**LATCH IMPLEMENTATIONS**

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Mellor-Crummey and Scott

![Diagram of CPU Latches]

- Base Latch
- CPU1 Latch
- CPU2 Latch
- CPU3
LATCH IMPLEMENTATIONS

Choice #4: Reader-Writer Locks
→ Allows for concurrent readers
→ Have to manage read/write queues to avoid starvation
→ Can be implemented on top of spinlocks

Latch

<table>
<thead>
<tr>
<th></th>
<th>read</th>
<th>write</th>
</tr>
</thead>
<tbody>
<tr>
<td>Latch</td>
<td>=0</td>
<td>=0</td>
</tr>
<tr>
<td>⌛</td>
<td>=0</td>
<td>=0</td>
</tr>
</tbody>
</table>
LATCH IMPLEMENTATIONS

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Latch

- read: 
  - \( L = 1 \)
  - \( R = 0 \)

- write:
  - \( L = 0 \)
  - \( R = 0 \)
LATCH IMPLEMENTATIONS

Choice #4: Reader-Writer Locks

→ Allows for concurrent readers
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LATCH CRABBING

Acquire and release latches on B+Tree nodes when traversing the data structure.

A thread can release latch on a parent node if its child node considered safe.
→ Any node that won’t split or merge when updated.
→ Not full (on insertion)
→ More than half-full (on deletion)
LATCH CRABBING

**Search:** Start at root and go down; repeatedly,
→ Acquire read (R) latch on child
→ Then unlock the parent node.

**Insert/Delete:** Start at root and go down, obtaining write (W) latches as needed.
Once child is locked, check if it is safe:
→ If child is safe, release all locks on ancestors.
EXAMPLE #1: SEARCH 23
EXAMPLE #1: SEARCH 23

R

A

20

B

10

D

6

E

12

23

F

38

G

35

44
We can release the latch on A as soon as we acquire the latch for C.
EXAMPLE #1: SEARCH 23

We can release the latch on A as soon as we acquire the latch for C.
EXAMPLE #1: SEARCH 23

We can release the latch on A as soon as we acquire the latch for C.
EXAMPLE #2: DELETE 44
EXAMPLE #2: DELETE 44

A

B

C

D

E

F

G

W

20

10

35

6

12

23

38

44
EXAMPLE #2: DELETE 44

We may need to coalesce C, so we can’t release the latch on A.
EXAMPLE #2: DELETE 44

We may need to coalesce C, so we can’t release the latch on A.

G will not merge with F, so we can release latches on A and C.
EXAMPLE #2: DELETE 44

We may need to coalesce C, so we can’t release the latch on A.

G will not merge with F, so we can release latches on A and C.
EXAMPLE #2: DELETE 44

We may need to coalesce C, so we can’t release the latch on A.

G will not merge with F, so we can release latches on A and C.
EXAMPLE #3: INSERT 40
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C has room if its child has to split, so we can release the latch on A.
EXAMPLE #3: INSERT 40

C has room if its child has to split, so we can release the latch on A.
EXAMPLE #3: INSERT 40

C has room if its child has to split, so we can release the latch on A.

G has to split, so we can’t release the latch on C.
EXAMPLE #3: INSERT 40

C has room if its child has to split, so we can release the latch on A.

G has to split, so we can’t release the latch on C.
EXAMPLE #3: INSERT 40

C has room if its child has to split, so we can release the latch on A.
G has to split, so we can’t release the latch on C.
EXAMPLE #3: INSERT 40

C has room if its child has to split, so we can release the latch on A.

G has to split, so we can’t release the latch on C.
OBSERVATION

What was the first step that the DBMS took in the two examples that updated the index?

Delete 44

Insert 40
BETTER LATCH CRABBING

Optimistically assume that the leaf is safe.
→ Take $R$ latches as you traverse the tree to reach it and verify.
→ If leaf is not safe, then do previous algorithm.

Also called *optimistic lock coupling*. 
EXAMPLE #4: DELETE 44
EXAMPLE #4: DELETE 44

We assume that C is safe, so we can release the latch on A.
EXAMPLE #4: DELETE 44

We assume that C is safe, so we can release the latch on A.
EXAMPLE #4: DELETE 44

We assume that C is safe, so we can release the latch on A.
Acquire an exclusive latch on G.
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EXAMPLE #4: DELETE 44

We assume that C is safe, so we can release the latch on A.

Acquire an exclusive latch on G.
OBSERVATION

Crabbing ensures that txns do not corrupt the internal data structure during modifications.

But because txns release latches on each node as soon as they are finished their operations, we cannot guarantee that phantoms do not occur…
PROBLEM SCENARIO #1

**Txn #1:** Check if 25 exists

Diagram: A is locked, and a request R is made to check if 25 exists.
PROBLEM SCENARIO #1

Txn #1: Check if 25 exists
PROBLEM SCENARIO #1

Txn #1: Check if 25 exists
PROBLEM SCENARIO #1

Txn #1: Check if 25 exists

Diagram:

A

B

C

D

E

F

G

20

10

35

6

12

23

38

44

Check if 25 exists
PROBLEM SCENARIO #1

Txn #1: Check if 25 exists
Txn #2: Insert 25
PROBLEM SCENARIO #1

**Txn #1:** Check if 25 exists

**Txn #2:** Insert 25
PROBLEM SCENARIO #1

Txn #1: Check if 25 exists
Txn #2: Insert 25
PROBLEM SCENARIO #1

Txn #1: Check if 25 exists
Txn #2: Insert 25
Txn #1: Insert 25
PROBLEM SCENARIO #1

Txn #1: Check if 25 exists
Txn #2: Insert 25
Txn #1: Insert 25
PROBLEM SCENARIO #2

A

B

C

D

E

F

G
PROBLEM SCENARIO #2

Txn #1: Scan [12, 23]
PROBLEM SCENARIO #2

Txn #1: Scan [12, 23]
**PROBLEM SCENARIO #2**

**Txn #1:** Scan [12, 23]

**Txn #2:** Insert 21
**PROBLEM SCENARIO #2**

**Txn #1:** Scan [12, 23]

**Txn #2:** Insert 21
**PROBLEM SCENARIO #2**

- **Txn #1:** Scan [12, 23]
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PROBLEM SCENARIO #2

 Txn #1: Scan [12, 23]
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PROBLEM SCENARIO #2

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PROBLEM SCENARIO #2

**Txn #1:** Scan [12, 23]
**Txn #2:** Insert 21
**Txn #1:** Scan [12, 23]
INDEX LOCKS

Need a way to protect the index’s logical contents from other txns to avoid phantoms.

Difference with index latches:
→ Locks are held for the entire duration of a txn.
→ Only acquired at the leaf nodes.
→ Not physically stored in index data structure.
INDEX LOCKS

Lock Table

txn1 X txn2 S  txn3 S

txn3 S  txn2 S  txn4 S

txn4 IX  txn6 X txn5 S
INDEX LOCKING SCHEMES

Predicate Locks
Key-Value Locks
Gap Locks
Key-Range Locks
Hierarchical Locking
PREDICATE LOCKS

Proposed locking scheme from System R.
→ Shared lock on the predicate in a WHERE clause of a SELECT query.
→ Exclusive lock on the predicate in a WHERE clause of any UPDATE, INSERT, or DELETE query.

Never implemented in any system.
PREDICATE LOCKS

SELECT SUM(balance) 
    FROM account 
    WHERE name = 'Biggie'

INSERT INTO account 
    (name, balance) 
    VALUES ('Biggie', 100);

Records in Table "account"
PREDICATE LOCKS

SELECT SUM(balance) FROM account WHERE name = 'Biggie'

INSERT INTO account (name, balance) VALUES ('Biggie', 100);

Records in Table "account"

name='Biggie'
PREDICATE LOCKS

SELECT SUM(balance) FROM account WHERE name = 'Biggie'

INSERT INTO account (name, balance) VALUES ('Biggie', 100);

Records in Table "account"
KEY-VALUE LOCKS

Locks that cover a single key value.
Need “virtual keys” for non-existent values.
GAP LOCKS

Each txn acquires a key-value lock on the single key that it wants to access. Then get a gap lock on the next key gap.

B+Tree Leaf Node

```
10 {Gap} 12 {Gap} 14 {Gap} 16
```
GAP LOCKS

Each txn acquires a key-value lock on the single key that it wants to access. Then get a gap lock on the next key gap.

*B+Tree Leaf Node*

| 10 {Gap} | 12 {Gap} | 14 {Gap} | 16 |

Gap (14, 16)
KEY- RANGE LOCKS

A txn takes locks on ranges in the key space.
→ Each range is from one key that appears in the relation, to the next that appears.
→ Define lock modes so conflict table will capture commutativity of the operations available.
KEY-RANGE LOCKS

Locks that cover a key value and the gap to the next key value in a single index.

→ Need “virtual keys” for artificial values (infinity)
KEY-RANGE LOCKS

Locks that cover a key value and the gap to the next key value in a single index.
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_B+Tree Leaf Node_

```
10 \{Gap\} 12 \{Gap\} 14 \{Gap\} 16
```

Prior Key (12, 14)
HIERARCHICAL LOCKING

Allow for a txn to hold wider key-range locks with different locking modes.
→ Reduces the number of visits to lock manager.

\[ \begin{array}{cccc}
10 & \text{Gap} & 12 & \text{Gap} \\
14 & \text{Gap} & 16 & \\
\end{array} \]
HIERARCHICAL LOCKING

Allow for a txn to hold wider key-range locks with different locking modes.
→ Reduces the number of visits to lock manager.

Tree Leaf Node

\[
\begin{array}{c}
\{10, 16\} \\
10 \{Gap\} \\
12 \{Gap\} \\
14 \{Gap\} \\
16
\end{array}
\]
HIERARCHICAL LOCKING

Allow for a txn to hold wider key-range locks with different locking modes.
→ Reduces the number of visits to lock manager.
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→ Reduces the number of visits to lock manager.
PARTING THOUGHTS

Hierarchical locking essentially provides predicate locking without complications.
→ Index locking occurs only in the leaf nodes.
→ Latching is to ensure consistent data structure.

Peloton currently does not support serializable isolation with range scans.
NEXT CLASS

Index Key Representation
Memory Allocation & Garbage Collection
T-Trees (1980s / TimesTen)
Bw-Tree (Hekaton)
Concurrent Skip Lists (MemSQL)